UC-Lab Center for Distribution System Cybersecurity

UCSB Presentation - Ramtin Pedarsani

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- The cyber-physical system is inherently geometrically distributed and has heterogeneous communications capability.



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• New energy management schemes need to be robust to network node and link failures.

- The goal is to maximize the welfare at time t for all generations and consumers.
- The ideal solution to the welfare maximization problem should have the following properties:
 - Manages the demand side and the generation side simultaneously.
 - Performs simple local optimization at each iteration and exchanges information with only neighbors.
 - Privacy of information is guaranteed. No information about the utility/cost functions should be disclosed.
 - Gonvergence to global optimum.
 - **(**) The algorithm should be scalable for large networks.

• We can formulate a constrained utility maximization problem:

Welfare maximization problem formulation

$$\max_{\mathbf{d},\mathbf{g}} \sum_{j\in\mathcal{J}} U_j(\mathbf{d}_j) - \sum_{v\in\mathcal{V}} C_v(\mathbf{g}_v)$$

s.t. $\mathbf{1}^T \mathbf{d} = \mathbf{1}^T \mathbf{g}$
 $d_{i,\min} \le d_i \le d_{i,\max}$
 $0 \le g_i \le g_{i,\max}$

• The t^{th} entry of the vectors of demand and generation correspond to period t (different times of the day).

- Next, for simplicity, we focus on one single period.
- Let $\Delta = (\sum_{j \in \mathcal{J}} d_j \sum_{v \in \mathcal{V}} g_v)$ be the power mismatch.

Dual Decomposition

$$J = \sum_{v \in \mathcal{V}} C_v(g_v) - \sum_{j \in \mathcal{J}} U_j(d_j) + p\Delta$$
$$d_i^{(k)} = \arg \min_{\substack{d_i, \min \leq d_i \leq d_i, \max \\ 0 \leq g_i \leq g_i, \max}} (p^{(k)}d_i - U_i(d_i))$$
$$g_i^{(k)} = \arg \min_{\substack{0 \leq g_i \leq g_i, \max \\ 0 \leq g_i \leq g_i, \max}} (C_i(d_i) - p^{(k)}g_i)$$
$$p^{(k+1)} = p^{(k)} + \eta\Delta^{(k)}$$

- $\bullet\,$ Global parameters are Δ and price p that are not private.
- In the consensus algorithm, each node can only share the global parameters with the neighbors.
- W is the weight matrix that also determines the communication graph.

Distributed Estimation of Global Information

$$\begin{split} i \in \mathcal{V}: \ \hat{\Delta}_{i}^{(k+1)} &= \hat{\Delta}_{i}^{(k)} + \sum_{j \in \mathcal{N}_{i}} w_{ij} (\hat{\Delta}_{j}^{(k)} - \hat{\Delta}_{i}^{(k)}) + g_{i}^{(k)} - g_{i}^{(k+1)} \\ i \in \mathcal{J}: \ \hat{\Delta}_{i}^{(k+1)} &= \hat{\Delta}_{i}^{(k)} + \sum_{j \in \mathcal{N}_{i}} w_{ij} (\hat{\Delta}_{j}^{(k)} - \hat{\Delta}_{i}^{(k)}) - d_{i}^{(k)} + d_{i}^{(k+1)} \\ i \in \mathcal{V} \cup \mathcal{J}: \ p_{i}^{(k+1)} &= p_{i}^{(k)} + \sum_{j \in \mathcal{N}_{i}} w_{ij} (p_{j}^{(k)} - p_{i}^{(k)}) + \eta \hat{\Delta}_{i}^{(k)} \end{split}$$

- We consider the attack that the adversary jams the communication links.
- The bandwidth on links will be significantly reduced. Can we solve the decentralized optimization problem?
- Key idea is to use quantization!

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$$\mathbf{x}_{i,t+1} = \sum_{\substack{j \in \mathcal{N}_i \\ \text{average of local} \\ \text{and neighboring models}}} w_{ij} \mathbf{x}_{j,t} \qquad \underbrace{-\alpha \nabla f_i(\mathbf{x}_{i,t})}_{\text{local gradient descent}}$$

[Nedić, Ozdaglar,'07]

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Theorem (Yuan, Ling, Yin '16)

Under A1,2,3, $\mathbf{x}_{i,t}$ geometrically converges to an $\mathcal{O}\left(\frac{\alpha}{1-\beta}\right)$ -neighborhood of the unique solution \mathbf{x}^* . $(1-\beta : \text{spectral gap of } W)$

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Related work in quantized setting: [Nedic et al., 2009], [Rabbat & Novak, 2005], [Li et al., 2016], [Zhang et al., 2019]

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No "EXACT" convergence!

Assumptions

$$\|
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$$|\nabla f_i(\mathbf{x}) - \nabla f_i(\mathbf{y})|| \le L ||\mathbf{x} - \mathbf{y}|| \qquad \forall \mathbf{x}, \mathbf{y}$$

A2 Local objectives f_i are μ -strongly convex:

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 & $W = W^{\top}$ & $W\mathbf{1} = \mathbf{1}$ & $\mathsf{null}(I - W) = \mathsf{span}(\mathbf{1})$

- Network $\mathcal G$ with n nodes, weight matrix $W \in \mathbb{R}^{n \times n}_+$
- At iteration *t*, node *i*:
 - sends $\mathbf{z}_{i,t} = Q(\mathbf{x}_{i,t})$ to neighbors $j \in \mathcal{N}_i$ and receives $\mathbf{z}_{j,t}$
 - updates

$$\mathbf{x}_{i,t+1} = \underbrace{(1 - \epsilon + \epsilon w_{ii})\mathbf{x}_{i,t}}_{\text{noiseless local model}} + \underbrace{\epsilon \sum_{j \in \mathcal{N}_i \setminus \{i\}} w_{ij} \mathbf{z}_{i,t}}_{j \in \mathcal{N}_i \setminus \{i\}} - \underbrace{-\alpha \epsilon \nabla f_i(\mathbf{x}_{i,t})}_{\text{local gradient descent}}$$

average of noisy neighboring models

- Network $\mathcal G$ with n nodes, weight matrix $W \in \mathbb{R}^{n \times n}_+$
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Theorem (QDGD with variance-bounded quantization)

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$$\epsilon=\frac{c_1}{T^{\frac{3}{4}(1-\delta)}}$$
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then

$$\mathbb{E}\left[\left\|\mathbf{x}_{i,T} - \mathbf{x}^*\right\|^2\right] \le \mathcal{O}\left(\frac{(1-\beta)^{-2} + n\sigma^2 \left\|W - W_D\right\|^2}{T^{\frac{1-\delta}{2}}}\right)$$

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A4 Random quantizer $Q(\cdot)$ is unbiased & variance-bounded:

 $\mathbb{E}[Q(\mathbf{x})|\mathbf{x}] = \mathbf{x} \quad \& \quad \mathbb{E}\left[\|Q(\mathbf{x}) - \mathbf{x}\|^2 |\mathbf{x}\right] \le \sigma^2 \qquad orall \mathbf{x}$

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• Solve an equivalent:

$$\begin{array}{ll} \min_{\mathbf{x} \in \mathbb{R}^{np}} & F(\mathbf{x}) = \sum_{i=1}^{n} f_i(\mathbf{x}_i) \\ \text{s.t.} & \mathbf{x}_1 = \dots = \mathbf{x}_n \end{array} \quad \text{where } \mathbf{x} = \begin{pmatrix} \mathbf{x}_1 \\ \vdots \\ \mathbf{x}_n \end{pmatrix} \in \mathbb{R}^{np}$$

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 where $\mathbf{W} = W \otimes I_p \in \mathbb{R}^{np \times np}$

Proof Sketch I

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• Penalty function:

Define
$$\forall \alpha$$
: $h_{\alpha}(\mathbf{x}) = \frac{1}{2} \mathbf{x}^{\top} (\mathbf{I} - \mathbf{W}) \mathbf{x} + \alpha F(\mathbf{x})$
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$$\mathbf{x}_{\alpha}^* = \operatorname*{argmin}_{\mathbf{x} \in \mathbb{R}^{np}} h_{\alpha}(\mathbf{x})$$









proposed update rule:
$$\mathbf{x}_{t+1} = \mathbf{x}_t - \epsilon \left(\underbrace{(\mathbf{I} - \mathbf{W}_D) \mathbf{x}_t + (\mathbf{W}_D - \mathbf{W}) \mathbf{z}_t + \alpha \nabla F(\mathbf{x}_t)}_{\overline{\nabla} h_\alpha(\mathbf{x}_t)} \& \mathbb{E}[\overline{\nabla} h_\alpha] = \nabla h_\alpha \right)$$



Stochastic gradient descent on penalty function $h_{\alpha}(\mathbf{x}_t)$:

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$$\Rightarrow \mathbb{E}\Big[\left\|\mathbf{x}_T - \mathbf{x}_{\alpha}^*\right\|^2\Big] \le \mathcal{O}\Big(\frac{c_1 n \sigma^2 \left\|W - W_D\right\|^2}{\mu c_2} \frac{1}{T^{(1-\delta)/2}}\Big) \quad \text{Lemma 1 } \checkmark$$





$$\mathbf{u}_{t+1} = \mathbf{u}_t - 1 \cdot \nabla h_\alpha(\mathbf{u}_t)$$



$$\begin{cases} \mathbf{u}_{t+1} = \mathbf{u}_t - 1 \cdot \nabla h_{\alpha}(\mathbf{u}_t) \Rightarrow \|\mathbf{u}_T - \mathbf{x}_{\alpha}^*\|^2 \le e^{-c_2 T^{(3+\delta)/4}} \|\mathbf{u}_0 - \mathbf{x}_{\alpha}^*\|^2 \end{cases}$$



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A4
$$\mathbb{E}[Q(\mathbf{x})|\mathbf{x}] = \mathbf{x}$$
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 $\mathsf{A4'} \ \mathbb{E}\left[Q(\mathbf{x})|\mathbf{x}\right] = \mathbf{x} \quad \& \quad \mathbb{E}\left[\|Q(\mathbf{x}) - \mathbf{x}\|^2 \, |\mathbf{x}\right] \leq \frac{\eta^2 \, \|\mathbf{x}\|^2}{\eta^2 \, \|\mathbf{x}\|^2} \qquad \text{for all } \mathbf{x}$

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Theorem

Under A1,2,3,4', the same rate is achieved.

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• Example: Low-precision Q.

 $Q_i^{\mathsf{LP}}(\mathbf{x}) = \|\mathbf{x}\| \cdot \mathsf{sign}(x_i) \cdot \xi_i(\mathbf{x}) \quad \& \quad \xi_i(\mathbf{x}) \text{ is a Bernoulli r.v. with parameter } \frac{|x_i|}{\|\mathbf{x}\|}$



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 $Q_i^{\rm LP}(\mathbf{x}) = \|\mathbf{x}\| \cdot {\rm sign}(x_i) \cdot \xi_i(\mathbf{x}) \quad \& \quad \xi_i(\mathbf{x}) \text{ is a Bernoulli r.v. with parameter}$

$$\frac{|x_i|}{\|\mathbf{x}\|}$$

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A4
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 & $\mathbb{E}[||Q(\mathbf{x}) - \mathbf{x}||^2 |\mathbf{x}] \le \sigma^2$ for all \mathbf{x}

A4' $\mathbb{E}[Q(\mathbf{x})|\mathbf{x}] = \mathbf{x}$ & $\mathbb{E}\left[\|Q(\mathbf{x}) - \mathbf{x}\|^2 |\mathbf{x}\right] \le \eta^2 \|\mathbf{x}\|^2$ for all \mathbf{x}

Theorem

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[1-bit SGD, Seide et al., 14], [QSGD, Alistarh, et al.,'17]

Numerical results: synthetic data

- Decentralized quadratic: $\min_{\mathbf{x}\in\mathbb{R}^p} f(\mathbf{x}) = \sum_{i=1}^n \frac{1}{2} \mathbf{x}^\top \mathbf{A}_i \mathbf{x} + \mathbf{b}_i^\top \mathbf{x}, \ p = 20$
- Network: Erdős-Rényi graph, n = 50 nodes, connectivity prob. $p_c = 0.35$
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s = 1	614.2	216.9	66.6
s = 10	11.69	678.2	3.96
$s^{*} = 50$	2.3	949.8	1.09
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communication cost to reach 0.01 of the optimal

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- We considered the attack that the adversary reduces communication bandwidth on the links, and proposed an exact decentralized gradient decent algorithm for quantized communications.
- Many interesting directions to continue:
 - Numerical study for IEEE 39-bus power network is ongoing.
 - best quantizer?
 - adversarial nodes?
 - Iink failures?
 - most resilient network topology?